

POSETS OF COPIES, EMBEDDING MONOIDS, AND INTERPRETABILITY OF RELATIONAL STRUCTURES

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Relational structures and complete Boolean algebras

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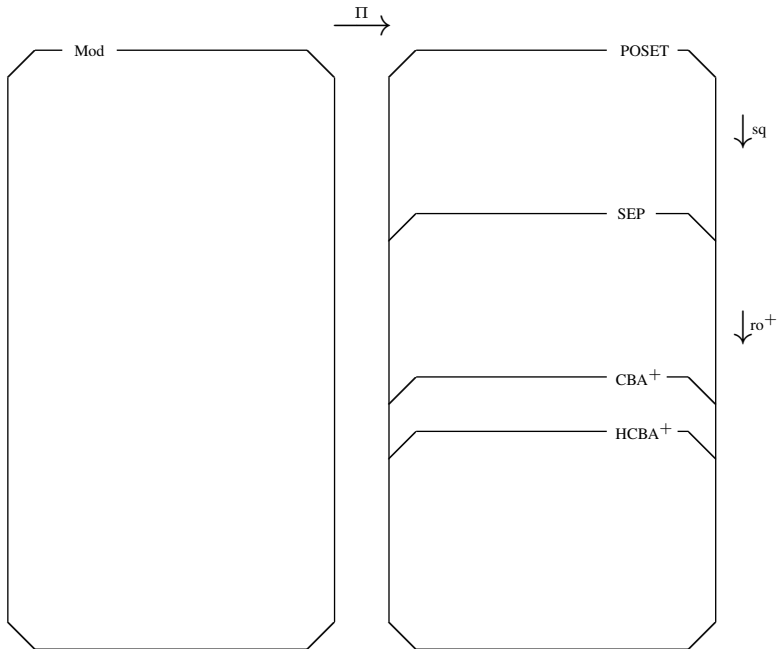
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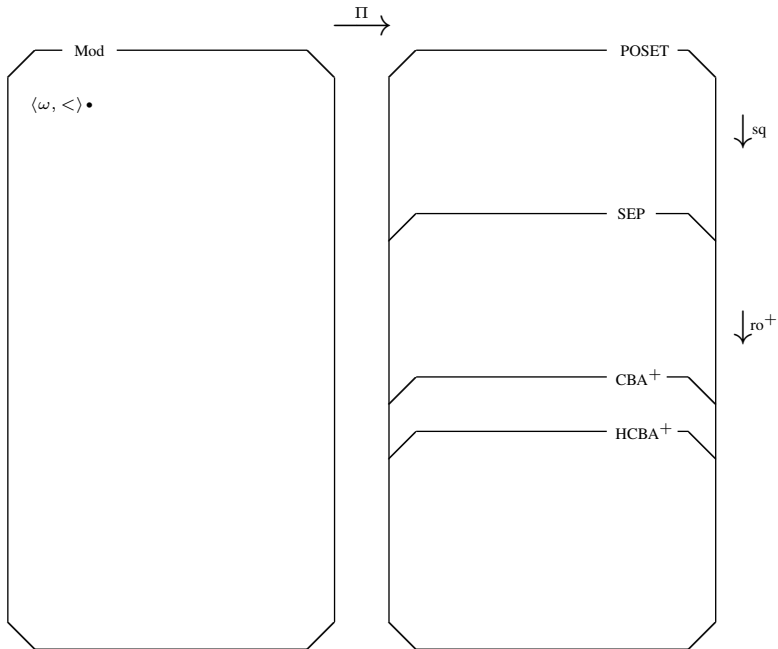
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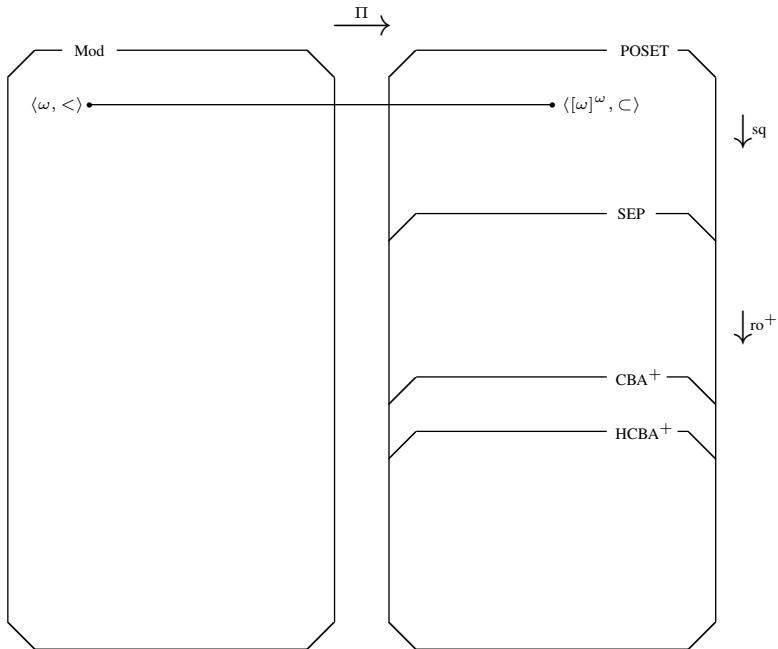
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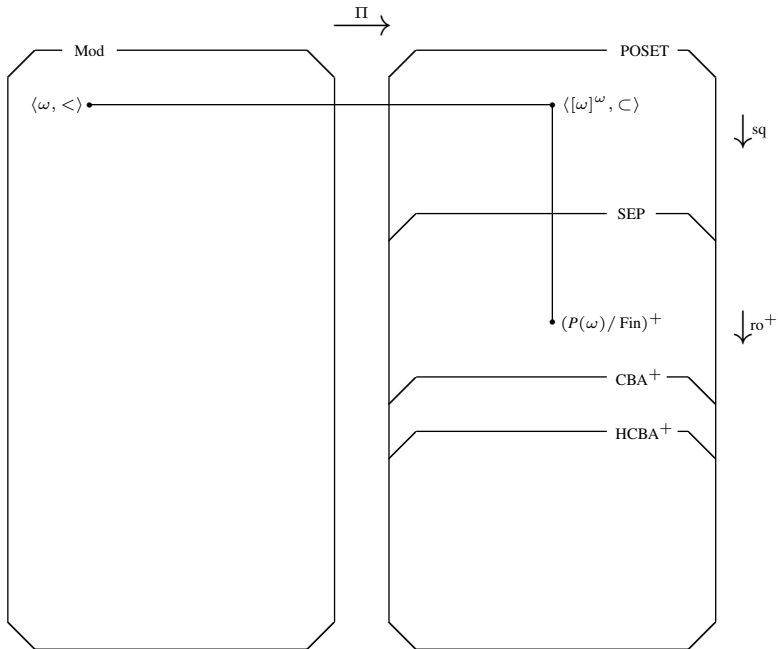
Theorem ([9])

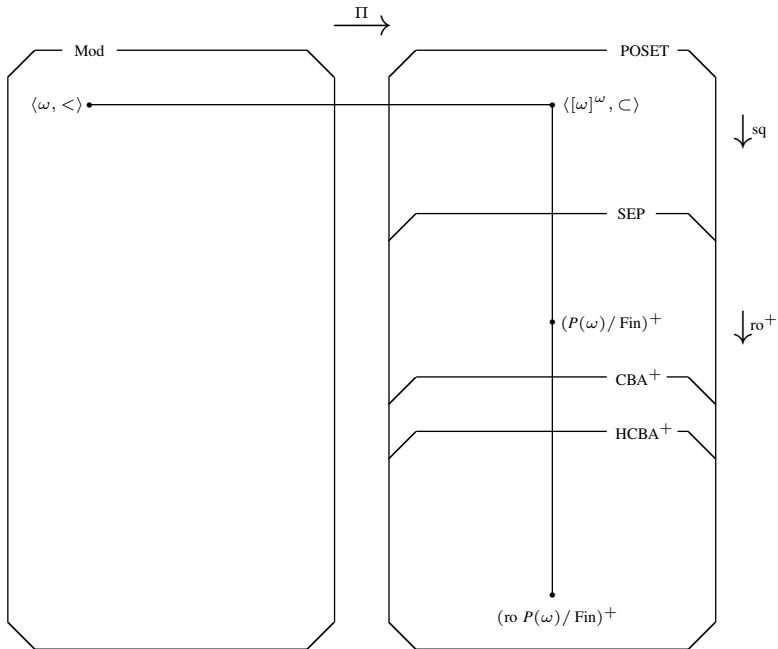
$\text{ro sq}\langle \mathbb{P}(\mathbb{X}), \subset \rangle$ is a homogeneous c. B. a.

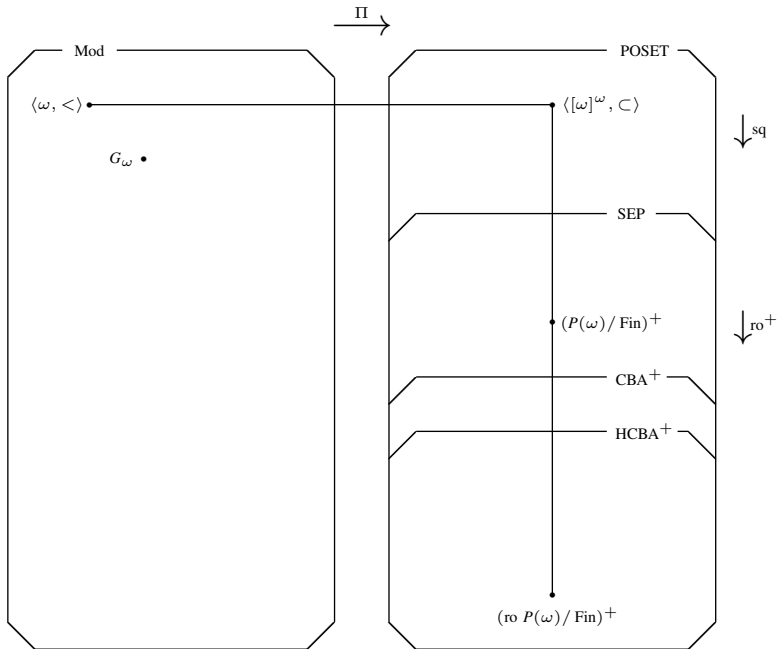


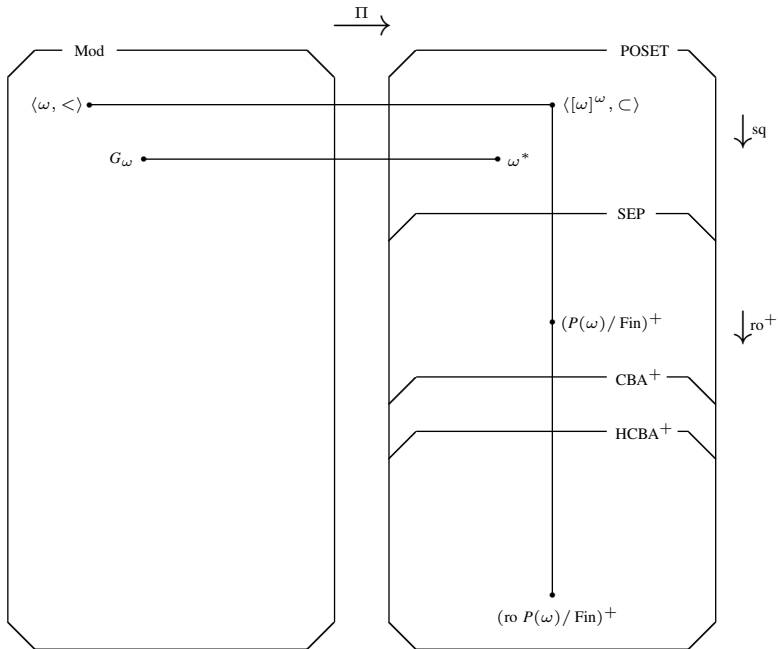


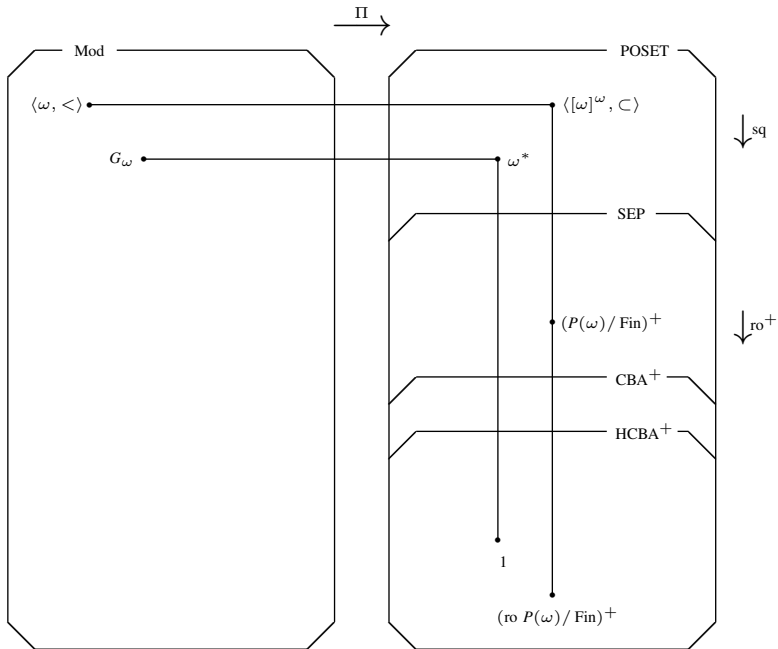


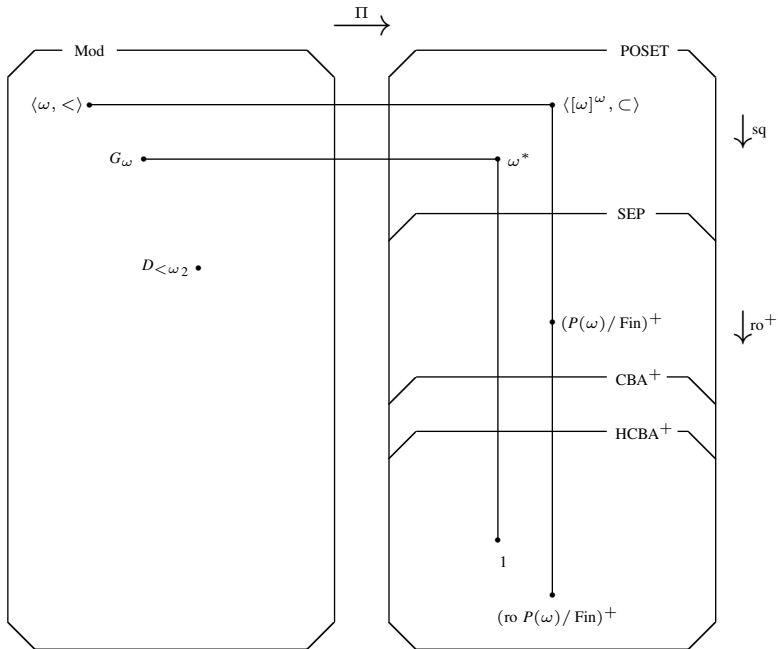


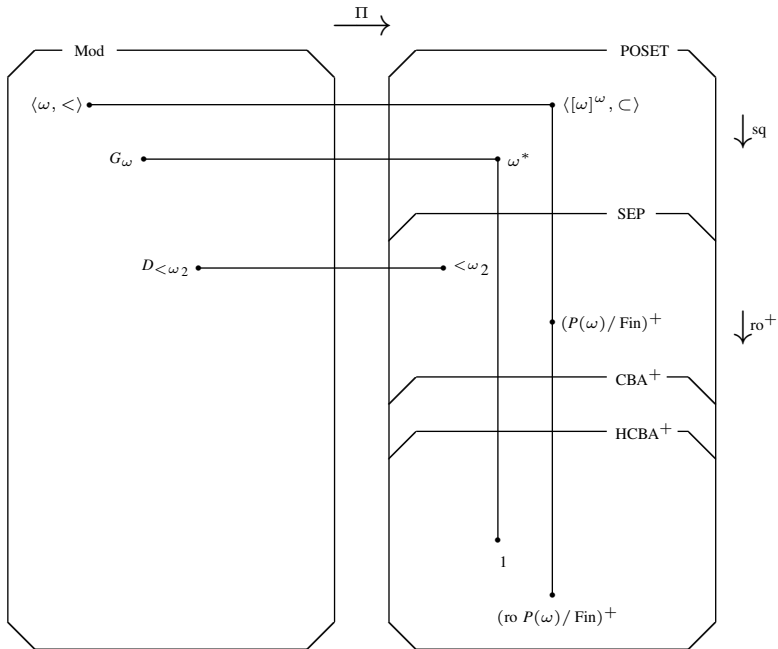


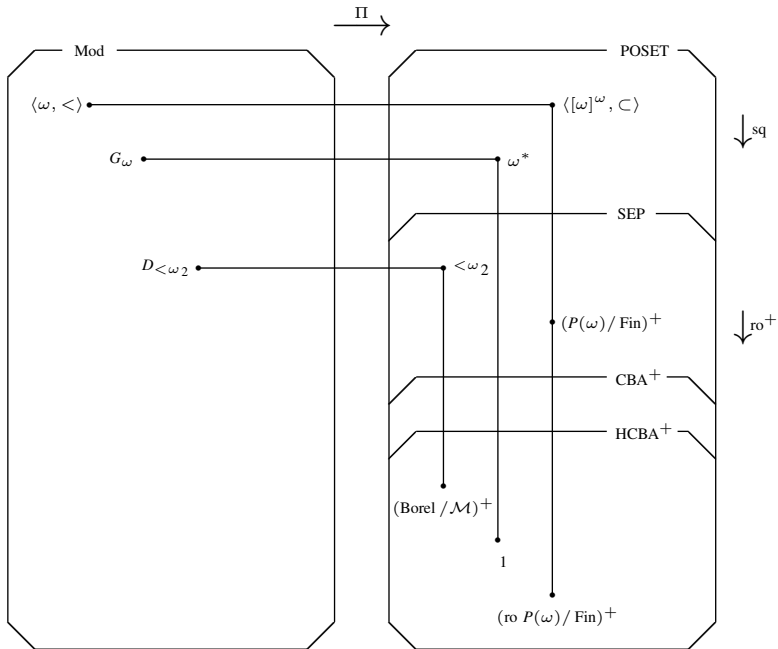


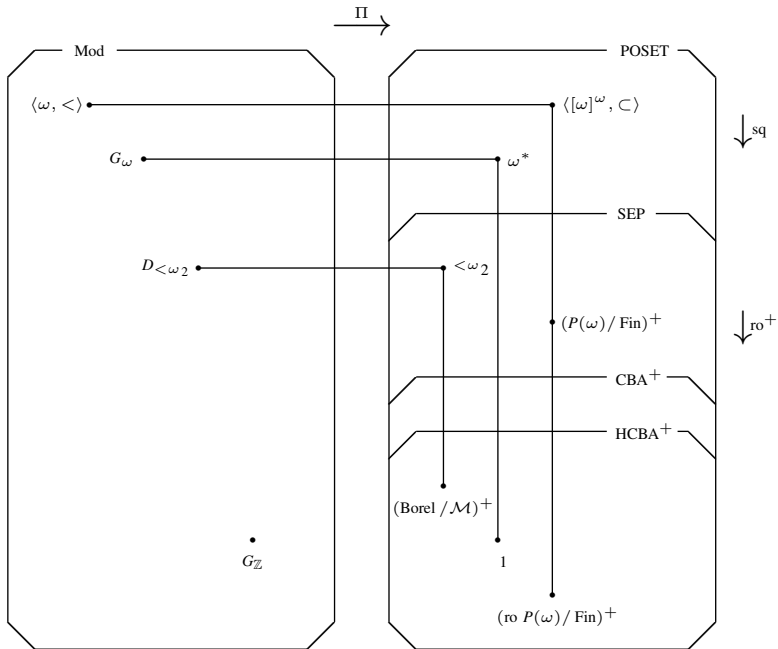


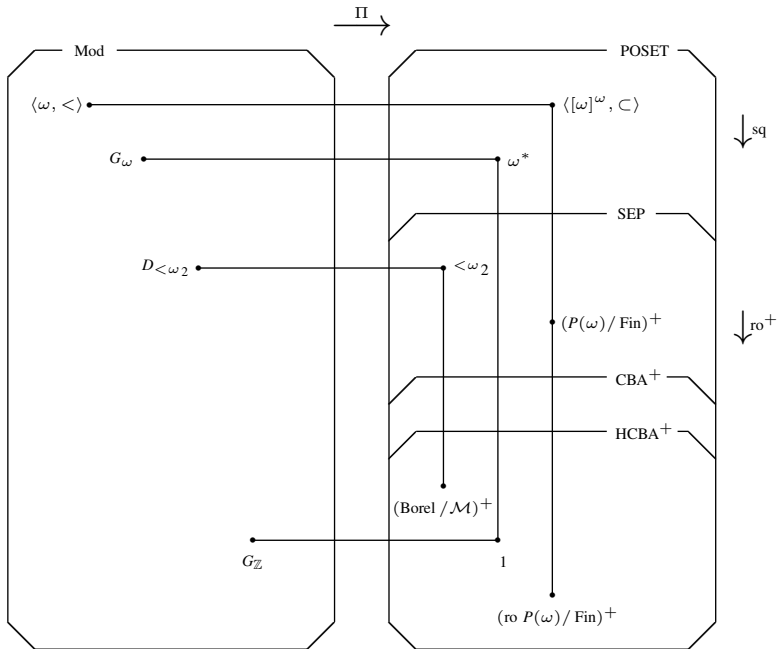




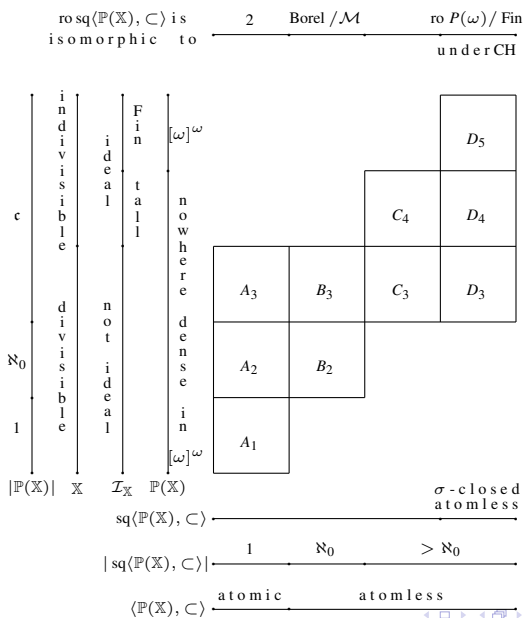




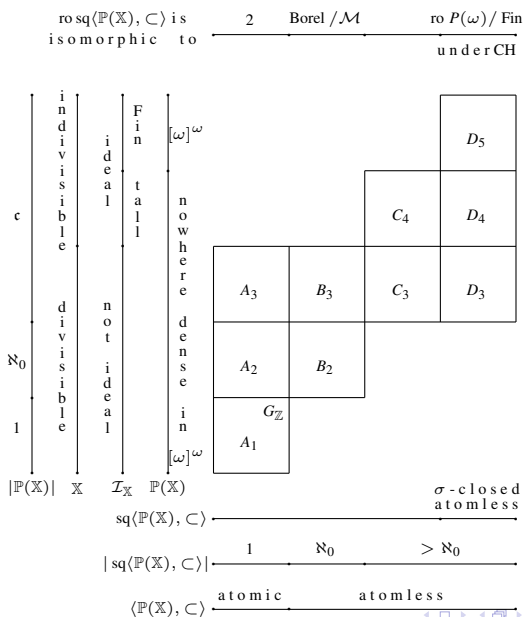




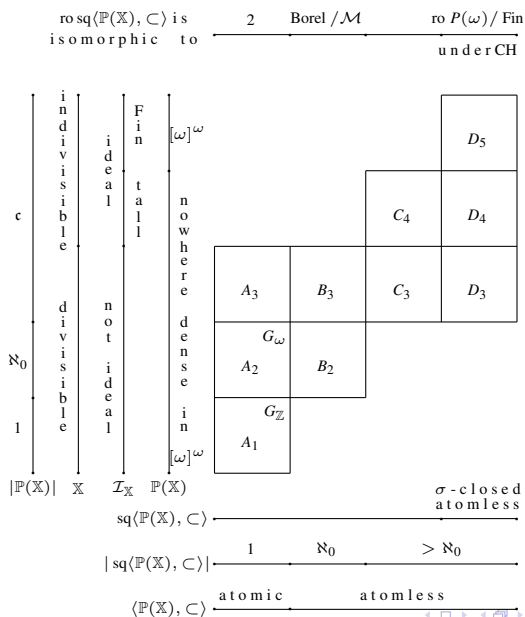
Countable binary structures [3]



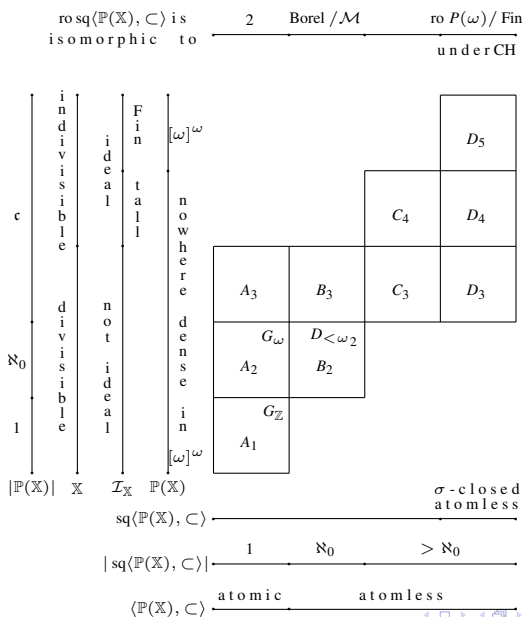
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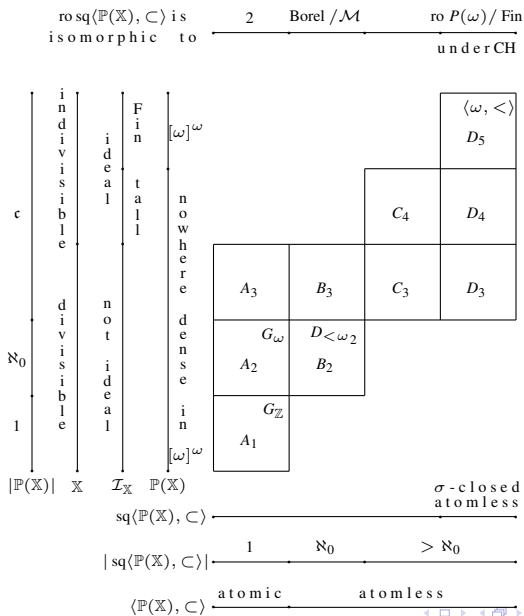
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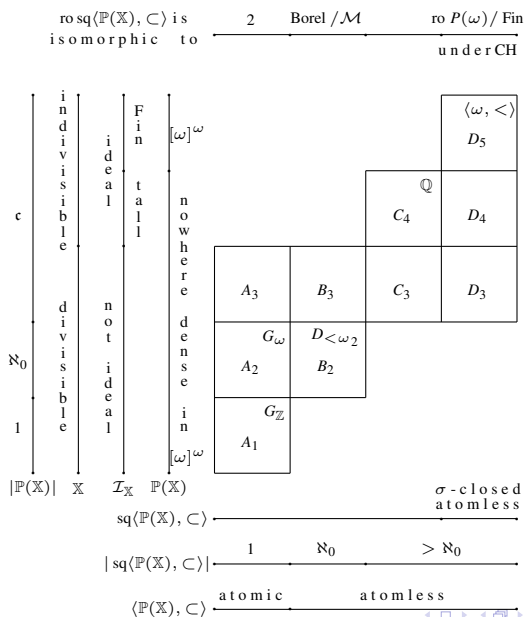
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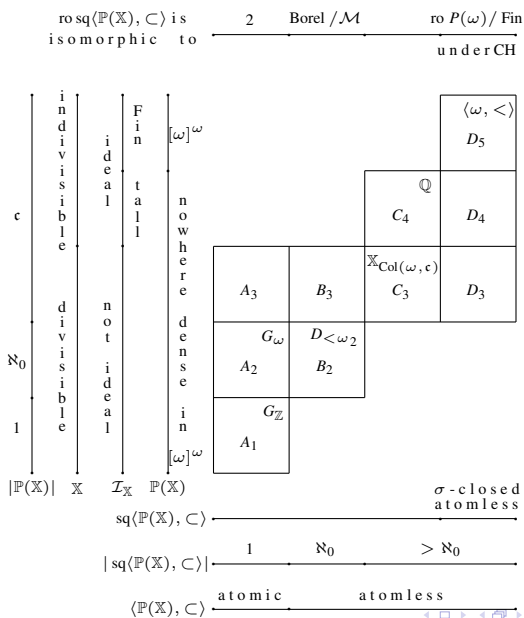
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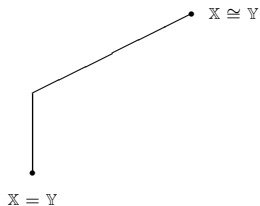
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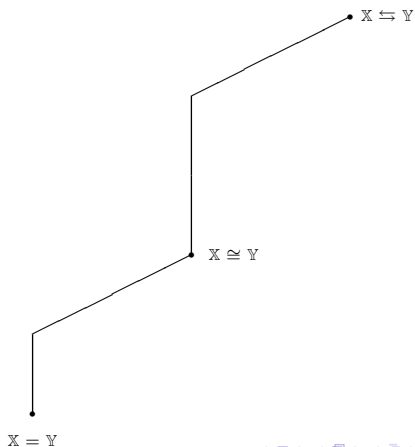
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$$X = Y$$

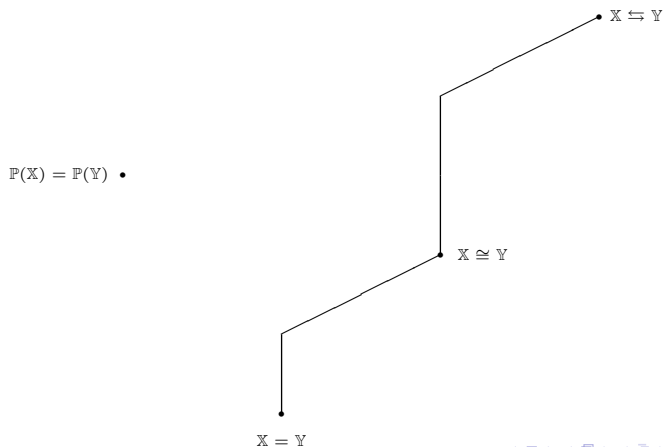
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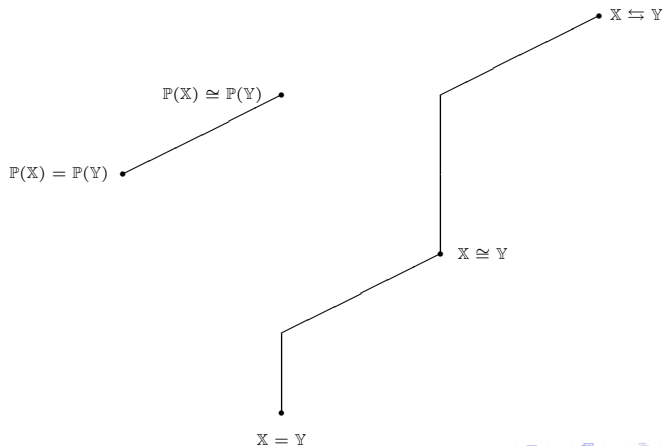
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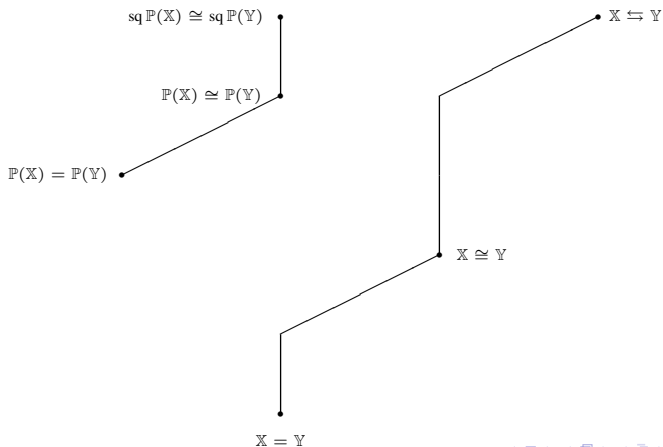
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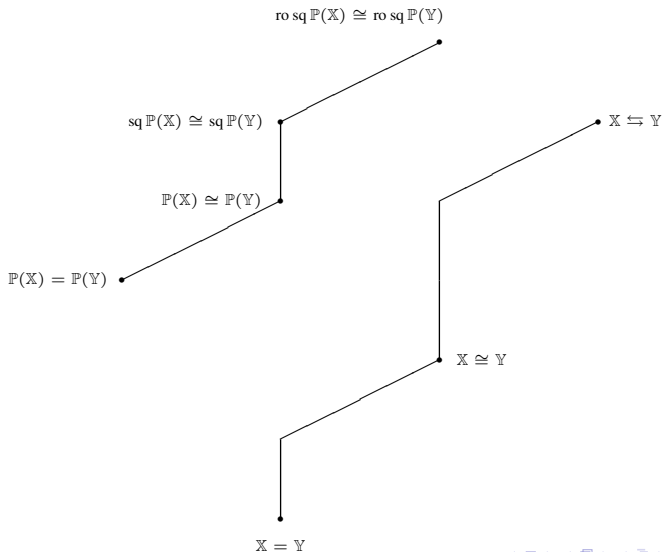
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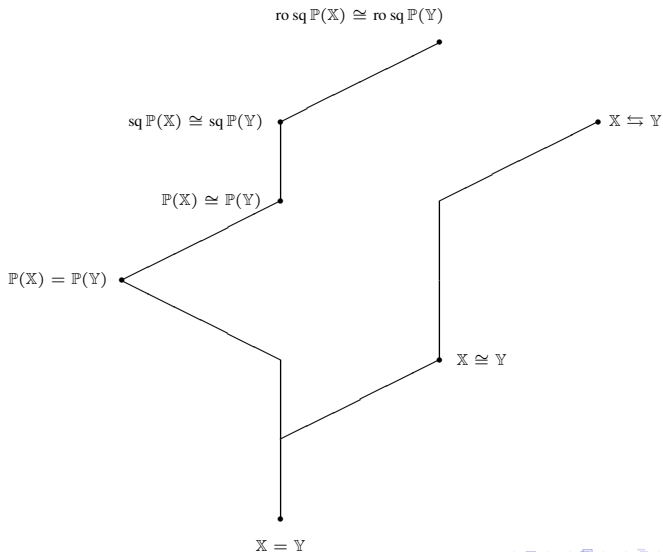
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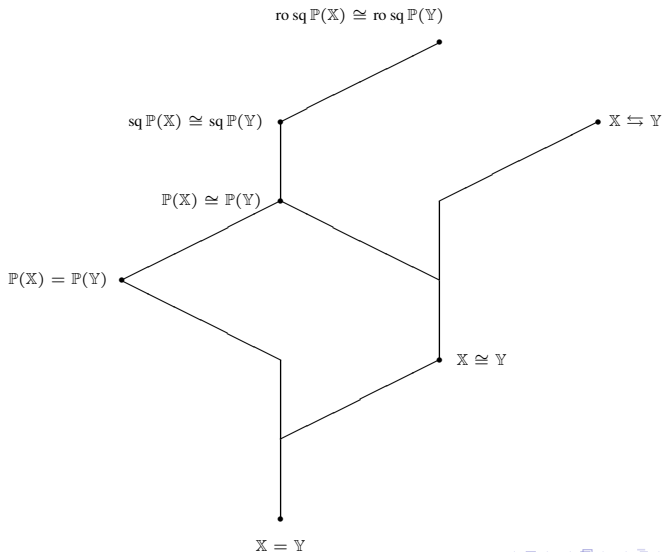
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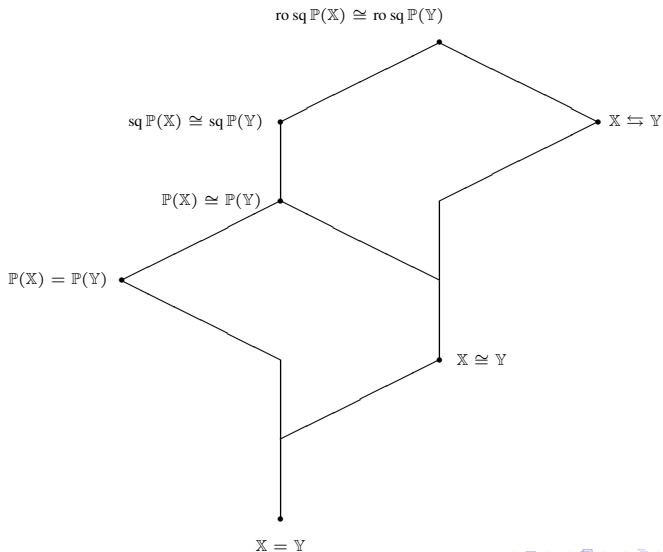
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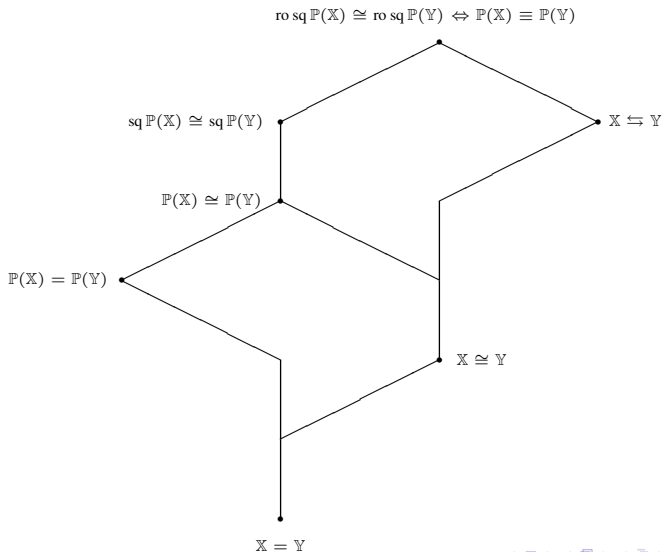
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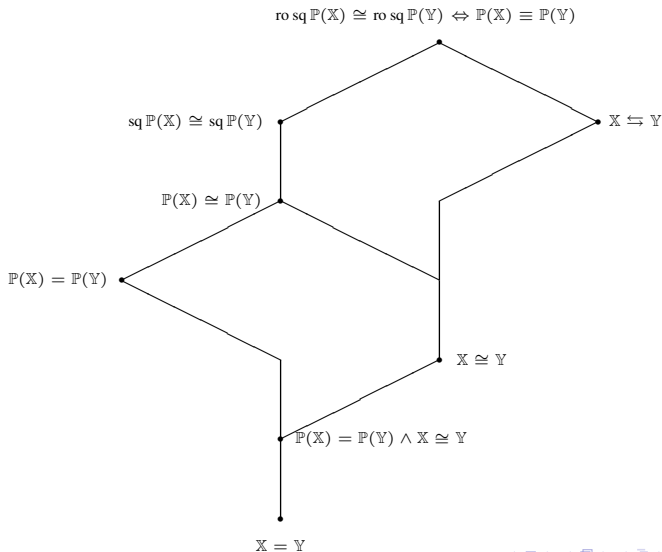
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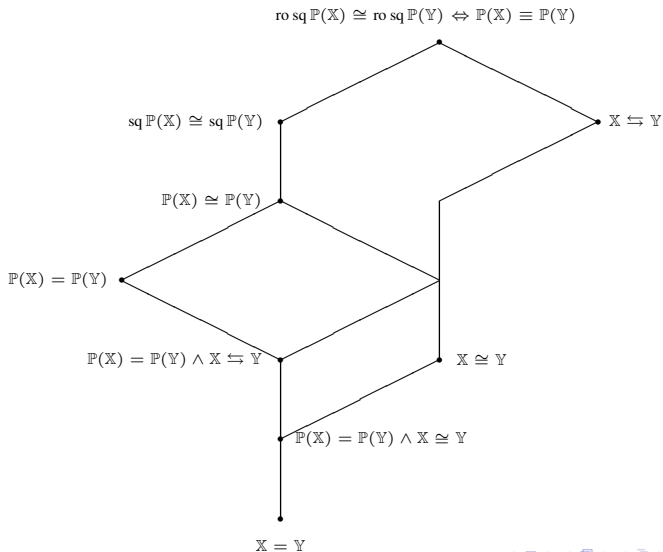
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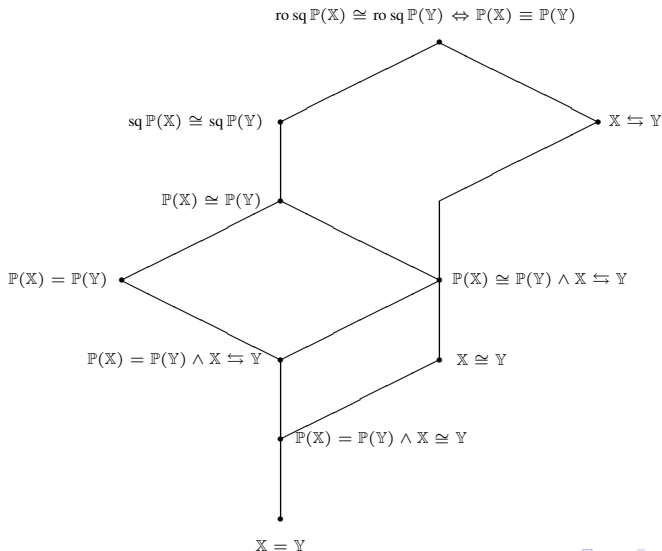
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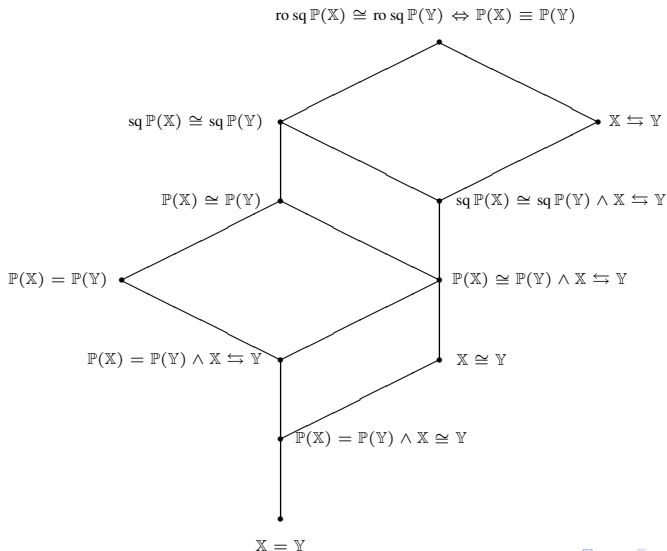
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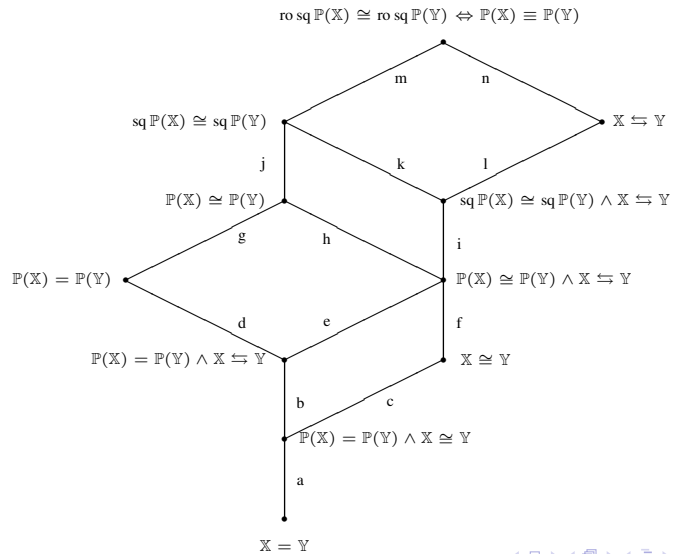
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- Under CH we have $\langle \mathbb{P}(\mathbb{X}), \subset \rangle \equiv (P(\omega) / \text{Fin})^+$.

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Theorem ([6])

If

$$\alpha = \omega^{\gamma_n+r_n} s_n + \cdots + \omega^{\gamma_0+r_0} s_0 + k$$

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$$\langle \mathbb{P}(\alpha), \subset \rangle \equiv \begin{cases} (P(\omega)/\text{Fin})^+ & \text{if } \alpha < \omega + \omega \\ (P(\omega)/\text{Fin})^+ * \pi & \text{if } \alpha \geq \omega + \omega \end{cases}$$

where $[\omega] \Vdash$ “ π is an ω_1 -closed, separative atomless forcing”.

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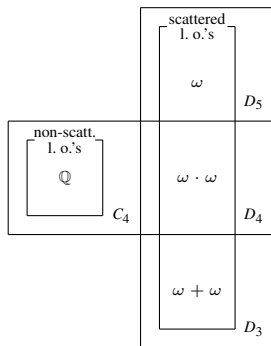
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- $1_{\mathbb{S}} \Vdash \pi \equiv (P(\omega)/\text{Fin})^+$, under CH or PFA.

Countable linear orders in the $A_1 - D_5$ diagram



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Theorem

If \mathbb{X} and \mathbb{Y} are relational structures, then

$$\mathbb{X} \cong \mathbb{Y} \Rightarrow \mathbf{Emb}(\mathbb{X}) \cong \mathbf{Emb}(\mathbb{Y}) \Rightarrow \langle \mathbb{P}(\mathbb{X}), \subset \rangle \cong \langle \mathbb{P}(\mathbb{Y}), \subset \rangle$$

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Since $\mathbb{P}((0, 1)_{\mathbb{Q}}, <) \not\cong \mathbb{P}([0, 1]_{\mathbb{Q}}, <)$ we have

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(Comment: but $\text{sq } \mathbb{P}((0, 1)_{\mathbb{Q}}, <) \cong \text{sq } \mathbb{P}([0, 1]_{\mathbb{Q}}, <)$)

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Proof. Using theorems of Grothendieck, Ore, and Dubreil.

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Theorem

\mathbb{X} and \mathbb{Y} are quantifier-free bi-interpretable $\Rightarrow \text{Emb}(\mathbb{X}) \cong \text{Emb}(\mathbb{Y})$.

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- Or $\text{sq } \mathbb{P}(\mathbb{X})$ is an atomless σ -closed poset (Column D) and, hence, $\mathbb{P}(\mathbb{X}) \equiv (P(\omega) / \text{Fin})^+$, under CH.

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