

On morphisms of lattice-valued formal contexts

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INVESTMENTS IN EDUCATION DEVELOPMENT

Outline

- 1 Introduction
- 2 Preliminaries on powerset operators
- 3 Categories of lattice-valued formal contexts
- 4 Properties of the categories of lattice-valued formal contexts
- 5 Conclusion

Formal Concept Analysis

- *Formal Concept Analysis (FCA)* has taken its origin as an attempt to restructure mathematics, e.g., lattice theory.
- Since then, FCA has been developed as a subfield of applied mathematics, based in mathematization of concept hierarchies.
- The aim of FCA is to support the rational communication of humans by mathematically developing appropriate conceptual structures, which can be logically activated.

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Formal contexts

One of the main building blocks of FCA provide *formal contexts*.

Definition 1

A *formal context* is a triple (G, M, I) , which comprises a set of objects G , a set of attributes M , and a binary incidence relation I between G and M , where $g I m$ means “object g has attribute m ”.

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Formal context morphisms

There exist at least three (different) ways of defining a morphism between two formal contexts (G_1, M_1, I_1) and (G_2, M_2, I_2) .

- ① The theory of FCA employs pairs of maps $G_1 \xrightarrow{\alpha} G_2, M_1 \xrightarrow{\beta} M_2$ such that $g I_1 m$ iff $\alpha(g) I_2 \beta(m)$ for every $g \in G_1, m \in M_1$.
- ② The theory of *Chu spaces* uses pairs of maps $G_1 \xrightarrow{\alpha} G_2, M_2 \xrightarrow{\beta} M_1$ such that $g I_1 \beta(m)$ iff $\alpha(g) I_2 m$ for every $g \in G_1, m \in M_2$.

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Formal context morphisms

- ③ The theory of *Galois connections* relies on the pairs of maps $\mathcal{P}(G_1) \xrightarrow{\alpha} \mathcal{P}(G_2)$, $\mathcal{P}(M_2) \xrightarrow{\beta} \mathcal{P}(M_1)$, where $\mathcal{P}(X)$ stands for the powerset of X , such that the diagrams

$$\begin{array}{ccc}
 \mathcal{P}(G_1) & \xrightarrow{\alpha} & \mathcal{P}(G_2) \\
 H_1 \downarrow & & \downarrow H_2 \\
 \mathcal{P}(M_1) & \xleftarrow{\beta} & \mathcal{P}(M_2)
 \end{array}
 \quad \text{and} \quad
 \begin{array}{ccc}
 \mathcal{P}(M_1) & \xleftarrow{\beta} & \mathcal{P}(M_2) \\
 K_1 \downarrow & & \downarrow K_2 \\
 \mathcal{P}(G_1) & \xrightarrow{\alpha} & \mathcal{P}(G_2)
 \end{array}$$

commute, where $H_j(S) = \{m \in M_j \mid s I_j m \text{ for every } s \in S\}$ and $K_j(T) = \{g \in G_j \mid g I_j t \text{ for every } t \in T\}$ (*Birkhoff operators*).

Lattice-valued formal contexts

- J. T. Denniston, A. Melton, and S. E. Rodabaugh compared the approaches of items (2) and (3) by considering their respective categories of *lattice-valued formal contexts* (in the sense of R. Bělohlávek) over a fixed commutative quantale Q , and constructing an embedding of each category into its counterparts.
- They finally arrived at the conclusion that the two viewpoints on formal context morphisms were not categorically isomorphic.

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Lattice-valued formal contexts

- This talk compares all three of the above-mentioned approaches to morphisms in the framework of lattice-valued formal contexts over a category of not necessarily commutative quantales.
- We construct a number of embeddings between their respective categories of formal contexts, showing that the approach of item (3) falls out of the FCA setting in the lattice-valued case.

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\vee -semilattices

Definition 2

CSLat(\vee) is the variety of *\vee -semilattices*, i.e., partially ordered sets (posets), which have arbitrary joins.

Every \vee -semilattice homomorphism $A_1 \xrightarrow{\varphi} A_2$ has the *upper adjoint map* $A_2 \xrightarrow{\varphi^\dagger} A_1$ given by $\varphi^\dagger(a_2) = \vee\{a_1 \in A_1 \mid \varphi(a_1) \leq a_2\}$.

∨-semilattices

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CSLat(∨) is the variety of *∨-semilattices*, i.e., partially ordered sets (posets), which have arbitrary joins.

Every ∨-semilattice homomorphism $A_1 \xrightarrow{\varphi} A_2$ has the *upper adjoint map* $A_2 \xrightarrow{\varphi^\dagger} A_1$ given by $\varphi^\dagger(a_2) = \bigvee \{a_1 \in A_1 \mid \varphi(a_1) \leq a_2\}$.

Quantaes

Definition 3

- ① **Quant** is the variety of *quantaes*, i.e., triples (Q, \vee, \otimes) , where
 - (Q, \vee) is a \vee -semilattice;
 - (Q, \otimes) is a semigroup;
 - \otimes distributes across \vee from both sides.
- ② **UQuant** is the variety of *unital quantaes*, i.e., quantaes Q , which have an element 1_Q such that $(Q, \otimes, 1_Q)$ is a monoid.

A quantale Q has two residuations, which are given by $q_1 \rightarrow_l q_2 = \vee \{q \in Q \mid q \otimes q_1 \leq q_2\}$ and $q_1 \rightarrow_r q_2 = \vee \{q \in Q \mid q_1 \otimes q \leq q_2\}$.

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Crisp forward powerset operator

Definition 4

Given a map $X_1 \xrightarrow{f} X_2$, the *forward powerset operator* w.r.t. f is the map $\mathcal{P}(X_1) \xrightarrow{f^\rightarrow} \mathcal{P}(X_2)$, which is defined by $f^\rightarrow(S) = \{f(s) \mid s \in S\}$.

Lattice-valued forward powerset operators I

Theorem 5

- Given a variety \mathbf{L} , which extends $\mathbf{CSLat}(\vee)$, every subcategory \mathbf{S} of \mathbf{L} provides a functor $\mathbf{Set} \times \mathbf{S} \xrightarrow{(-)^{\rightarrow}} \mathbf{CSLat}(\vee)$, which is defined by $((X_1, L_1) \xrightarrow{(f, \varphi)} (X_2, L_2))^{\rightarrow} = L_1^{X_1} \xrightarrow{(f, \varphi)^{\rightarrow}} L_2^{X_2}$, where $((f, \varphi)^{\rightarrow}(\alpha))(x_2) = \varphi(\bigvee_{f(x_1)=x_2} \alpha(x_1))$.
- Let \mathbf{L} be a variety, which extends $\mathbf{CSLat}(\vee)$, and let \mathbf{S} be a subcategory of \mathbf{L}^{op} such that for every \mathbf{S} -morphism $L_1 \xrightarrow{\varphi} L_2$, the map $L_1 \xrightarrow{\varphi^{op\uparrow}} L_2$ is \vee -preserving. Then there exists a functor $\mathbf{Set} \times \mathbf{S} \xrightarrow{(-)^{\uparrow\rightarrow}} \mathbf{CSLat}(\vee)$ defined by $((X_1, L_1) \xrightarrow{(f, \varphi)} (X_2, L_2))^{\uparrow\rightarrow} = L_1^{X_1} \xrightarrow{(f, \varphi)^{\uparrow\rightarrow}} L_2^{X_2}$, where $((f, \varphi)^{\uparrow\rightarrow}(\alpha))(x_2) = \varphi^{op\uparrow}(\bigvee_{f(x_1)=x_2} \alpha(x_1))$.

Lattice-valued forward powerset operators II

Theorem 6

- Given a variety \mathbf{L} , which extends $\mathbf{CSLat}(\vee)$, every subcategory \mathbf{S} of \mathbf{L}^{op} provides a functor $\mathbf{Set}^{op} \times \mathbf{S} \xrightarrow{(-)^{\rightarrow o}} (\mathbf{CSLat}(\vee))^{op}$ with $((X_1, L_1) \xrightarrow{(f, \varphi)} (X_2, L_2))^{\rightarrow o} = L_1^{X_1} \xrightarrow{((f, \varphi)^{\rightarrow o})^{op}} L_2^{X_2}$, where $((f, \varphi)^{\rightarrow o}(\alpha))(x_1) = \varphi^{op}(\vee_{f^{op}(x_2)=x_1} \alpha(x_2))$.
- Let \mathbf{L} be a variety, which extends $\mathbf{CSLat}(\vee)$, and let \mathbf{S} be a subcategory of \mathbf{L} such that for every \mathbf{S} -morphism $L_1 \xrightarrow{\varphi} L_2$, the map $L_2 \xrightarrow{\varphi^\dagger} L_1$ is \vee -preserving. Then there exists a functor $\mathbf{Set}^{op} \times \mathbf{S} \xrightarrow{(-)^{\dagger \rightarrow o}} (\mathbf{CSLat}(\vee))^{op}$ defined by $((X_1, L_1) \xrightarrow{(f, \varphi)} (X_2, L_2))^{\dagger \rightarrow o} = L_1^{X_1} \xrightarrow{((f, \varphi)^{\dagger \rightarrow o})^{op}} L_2^{X_2}$, where $((f, \varphi)^{\dagger \rightarrow o}(\alpha))(x_1) = \varphi^\dagger(\vee_{f^{op}(x_2)=x_1} \alpha(x_2))$.

Galois connections

Definition 7

A tuple $((X_1, \leq), f, g, (X_2, \leq))$ is an *order-reversing Galois connection* provided that $(X_1, \leq), (X_2, \leq)$ are posets, and $X_1 \begin{matrix} \xrightarrow{f} \\ \xleftarrow{g} \end{matrix} X_2$ are maps with $x_1 \leq g(x_2)$ iff $x_2 \leq f(x_1)$ for every $x_1 \in X_1, x_2 \in X_2$.

Formal contexts as Chu spaces

Definition 8

Let \mathbf{L} be a variety, which extends \mathbf{Quant} , and let \mathbf{S} be a subcategory of \mathbf{L}^{op} . $\mathbf{S}\text{-FC}^C$ is the category, which comprises the following data.

Objects: tuples $\mathcal{K} = (G, M, L, I)$ (*lattice-valued formal contexts*), where G is the set of context *objects*, M is the set of context *attributes*, L is an \mathbf{S} -object, and $G \times M \xrightarrow{I} L$ is a map, which is called the context *incidence relation*.

Morphisms: $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ (*lattice-valued formal context morphisms*) are triples $(G_1, M_1, L_1) \xrightarrow{f=(\alpha, \beta, \varphi)} (G_2, M_2, L_2)$ in $\mathbf{Set} \times \mathbf{Set}^{op} \times \mathbf{S}$ with $I_1(g, \beta^{op}(m)) = \varphi^{op} \circ I_2(\alpha(g), m)$ for every $g \in G_1, m \in M_2$.

Modified formal contexts as Chu spaces

Definition 9

Let \mathbf{L} be a variety, which extends **Quant**, and let \mathbf{S} be a subcategory of \mathbf{L} . $\mathbf{S}\text{-FC}_m^C$ is the category, which comprises the following data.

Objects: (lattice-valued) formal contexts.

Morphisms: $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ are triples $(G_1, M_1, L_1) \xrightarrow{f=(\alpha, \beta, \varphi)} (G_2, M_2, L_2)$ in $\mathbf{Set} \times \mathbf{Set}^{op} \times \mathbf{S}$ with $\varphi \circ l_1(g, \beta^{op}(m)) = l_2(\alpha(g), m)$ for every $g \in G_1, m \in M_2$.

Formal contexts of B. Ganter and R. Wille

Definition 10

Let \mathbf{L} be a variety, which extends **Quant**, and let \mathbf{S} be a subcategory of \mathbf{L}^{op} . $\mathbf{S}\text{-FC}^{GW}$ is the category, which comprises the following data.

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Definition 11

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Objects: (lattice-valued) formal contexts.

Morphisms: $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ are triples $(G_1, M_1, L_1) \xrightarrow{f=(\alpha, \beta, \varphi)} (G_2, M_2, L_2)$ in $\mathbf{Set} \times \mathbf{Set} \times \mathbf{S}$ with $\varphi \circ l_1(g, m) = l_2(\alpha(g), \beta(m))$ for every $g \in G_1, m \in M_1$.

Lattice-valued Birkhoff operators

Definition 12

Every lattice-valued formal context \mathcal{K} provides the following (*lattice-valued*) *Birkhoff operators*:

- ① $L^G \xrightarrow{H} L^M$ given by $(H(s))(m) = \bigwedge_{g \in G} (s(g) \rightarrow_l I(g, m))$;
- ② $L^M \xrightarrow{K} L^G$ given by $(K(t))(g) = \bigwedge_{m \in M} (t(m) \rightarrow_r I(g, m))$.

Theorem 13

For every lattice-valued context \mathcal{K} , (L^G, H, K, L^M) is an order-reversing Galois connection.

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For every lattice-valued context \mathcal{K} , (L^G, H, K, L^M) is an order-reversing Galois connection.

Crisp Birkhoff operators

Example 14

Every crisp context \mathcal{K} provides the maps

$$\textcircled{1} \mathcal{P}(G) \xrightarrow{H} \mathcal{P}(M), H(S) = \{m \in M \mid s I m \text{ for every } s \in S\};$$

$$\textcircled{2} \mathcal{P}(M) \xrightarrow{K} \mathcal{P}(G), K(T) = \{g \in G \mid g I t \text{ for every } t \in T\};$$

which are the classical Birkhoff operators of a binary relation.

Formal contexts of J. T. Denniston *et al.*

Definition 15

Given a variety \mathbf{L} , which extends \mathbf{Quant} , and a subcategory \mathbf{S} of \mathbf{L} , $\mathbf{S}\text{-FC}^{DMR}$ is the category, concrete over the product category $\mathbf{Set} \times \mathbf{Set}^{op}$, which comprises the following data.

Objects: lattice-valued formal contexts \mathcal{K} with L an object of \mathbf{S} .

Morphisms: $\mathcal{K}_1 \xrightarrow{f=(\alpha,\beta)} \mathcal{K}_2$ are $\mathbf{Set} \times \mathbf{Set}^{op}$ -morphisms $(L_1^{G_1}, L_1^{M_1}) \xrightarrow{(\alpha,\beta)} (L_2^{G_2}, L_2^{M_2})$, making the next diagrams commute

$$\begin{array}{ccc}
 L_1^{G_1} & \xrightarrow{\alpha} & L_2^{G_2} \\
 H_1 \downarrow & & \downarrow H_2 \\
 L_1^{M_1} & \xleftarrow{\beta^{op}} & L_2^{M_2}
 \end{array}$$

$$\begin{array}{ccc}
 L_1^{M_1} & \xleftarrow{\beta^{op}} & L_2^{M_2} \\
 K_1 \downarrow & & \downarrow K_2 \\
 L_1^{G_1} & \xrightarrow{\alpha} & L_2^{G_2}
 \end{array}$$

Relations versus Birkhoff operators

- There is a one-to-one correspondence between relations $I \subseteq G \times M$ and order-reversing Galois connections on $(\mathcal{P}(G), \mathcal{P}(M))$.
- What about the lattice-valued case?

Definition 16

Given a \vee -semilattice L and a set X , every $S \subseteq X$ and every $a \in L$ provide the map $X \xrightarrow{\chi_S^a} L$, which is defined by

$$\chi_S^a(x) = \begin{cases} a, & x \in S \\ \perp_L, & \text{otherwise.} \end{cases}$$

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Lattice-valued relations versus Birkhoff operators

Theorem 17

Let G, M be sets and let L be a unital quantale. For every order-reversing Galois connection $(L^G, \alpha, \beta, L^M)$, equivalent are:

- ① There exists a map $G \times M \xrightarrow{!} L$ such that $\alpha = H$ and $\beta = K$.
- ② For every $g \in G, m \in M, a \in L$, it follows that
 - (a) $(\alpha(\chi_{\{g\}}^{1L}))(m) = (\beta(\chi_{\{m\}}^{1L}))(g)$;
 - (b) $(\alpha(\underline{a} \otimes \chi_{\{g\}}^{1L}))(m) = a \rightarrow_l (\alpha(\chi_{\{g\}}^{1L}))(m)$;
 - (c) $(\beta(\chi_{\{m\}}^{1L} \otimes \underline{a}))(g) = a \rightarrow_r (\beta(\chi_{\{m\}}^{1L}))(g)$.
- ③ For every $g \in G, m \in M, a \in L$, it follows that
 - (a) $(\alpha(\underline{a} \otimes \chi_{\{g\}}^{1L}))(m) = a \rightarrow_l (\beta(\chi_{\{m\}}^{1L}))(g)$;
 - (b) $(\beta(\chi_{\{m\}}^{1L} \otimes \underline{a}))(g) = a \rightarrow_r (\alpha(\chi_{\{g\}}^{1L}))(m)$.

Consequences

Every map $G \times M \xrightarrow{I} L$ gives rise to an order-reversing Galois connection, but the converse way needs additional requirements.

Counterexample

Let L be the unit interval $\mathbb{I} = ([0, 1], \vee, \wedge, 1)$, and let both G and M be singletons. One can assume that both \mathbb{I}^G and \mathbb{I}^M is \mathbb{I} . The order-reversing involution map $\mathbb{I} \xrightarrow{\alpha} \mathbb{I}$, $\alpha(a) = 1 - a$ is a part of the order-reversing Galois connection $(\mathbb{I}, \alpha, \alpha, \mathbb{I})$. The condition of, e.g., Theorem 17(3)(a) gives $\alpha(a) = a \rightarrow \alpha(1)$ for every $a \in \mathbb{I}$. However, for $a = \frac{1}{2}$, one obtains that $\alpha(\frac{1}{2}) = \frac{1}{2} \neq 0 = \frac{1}{2} \rightarrow 0 = \frac{1}{2} \rightarrow \alpha(1)$.

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From $\mathbf{S}\text{-FC}^C$ to $\mathbf{S}\text{-FC}^{DMR}$

Definition 18

- $\mathbf{S}\text{-FC}_*^C$ is a subcategory of $\mathbf{S}\text{-FC}^C$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have surjective maps $G_1 \xrightarrow{\alpha} G_2$, $M_2 \xrightarrow{\beta^{op}} M_1$, and an \mathbf{S} -isomorphism $L_1 \xrightarrow{\varphi} L_2$.
- Let \mathbf{L} extend \mathbf{UQuant} . $\mathbf{S}\text{-FC}_{**}^C$ (resp. $\mathbf{S}\text{-FC}_{*\bullet}^C$) is a full subcategory of $\mathbf{S}\text{-FC}_*^C$, whose objects $\mathcal{K} = (G, M, L, I)$ have non-empty G (resp. M) and, moreover, $\perp_L \neq \perp_{L'}$.

Theorem 19

There exists a functor $\mathbf{S}\text{-FC}_*^C \xrightarrow{H_{CD}} \mathbf{S}\text{-FC}^{DMR}$, which is given by $H_{CD}(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = \mathcal{K}_1 \xrightarrow{((\alpha, \varphi)^{\perp \rightarrow \circ}, ((\beta, \varphi)^{\rightarrow \circ})^{op})} \mathcal{K}_2$. Its restriction to $\mathbf{S}\text{-FC}_{**}^C$ (resp. $\mathbf{S}\text{-FC}_{*\bullet}^C$) is a (non-full) embedding.

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From $\mathbf{S}\text{-FC}_m^C$ to $\mathbf{S}\text{-FC}^{DMR}$

Definition 20

- $\mathbf{S}\text{-FC}_{m^*}^C$ is a subcategory of $\mathbf{S}\text{-FC}_m^C$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have surjective maps $G_1 \xrightarrow{\alpha} G_2$, $M_2 \xrightarrow{\beta^{op}} M_1$, and an \mathbf{S} -isomorphism $L_1 \xrightarrow{\varphi} L_2$.
- Let \mathbf{L} extend \mathbf{UQuant} . $\mathbf{S}\text{-FC}_{m^{**}}^C$ (resp. $\mathbf{S}\text{-FC}_{m^{*\bullet}}^C$) is a full subcategory of $\mathbf{S}\text{-FC}_{m^*}^C$, whose objects $\mathcal{K} = (G, M, L, I)$ have non-empty G (resp. M) and, moreover, $1_L \neq \perp_L$.

Theorem 21

There exists a functor $\mathbf{S}\text{-FC}_{m^*}^C \xrightarrow{H_{CmD}} \mathbf{S}\text{-FC}^{DMR}$, which is given by $H_{CmD}(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = \mathcal{K}_1 \xrightarrow{((\alpha, \varphi) \rightarrow, ((\beta, \varphi)^{\dashv} \dashrightarrow \circ)^{op})} \mathcal{K}_2$. Its restriction to $\mathbf{S}\text{-FC}_{m^{**}}^C$ (resp. $\mathbf{S}\text{-FC}_{m^{*\bullet}}^C$) is a (non-full) embedding.

From $\mathbf{S}\text{-FC}_m^C$ to $\mathbf{S}\text{-FC}^{DMR}$

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Formal concepts, protoconcepts, and preconcepts

Definition 22

Let \mathcal{K} be a lattice-valued formal context, and let $s \in L^G$, $t \in L^M$. The pair (s, t) is called a

- *(lattice-valued) formal concept of \mathcal{K}* provided that $H(s) = t$ and $K(t) = s$;
- *(lattice-valued) formal protoconcept of \mathcal{K}* provided that $K \circ H(s) = K(t)$ (equivalently, $H \circ K(t) = H(s)$);
- *(lattice-valued) formal preconcept of \mathcal{K}* provided that $s \leq K(t)$ (equivalently, $t \leq H(s)$).

From $\mathbf{S}\text{-FC}^{\text{DMR}}$ to $\mathbf{S}\text{-FC}^{\text{C}}$

Definition 23

- Given an \mathbf{L} -algebra L , $L\text{-FC}_i^{\text{DMR}}$ is a subcategory of $L\text{-FC}^{\text{DMR}}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have injective maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{\text{op}}} L^{M_1}$.
- An \mathbf{L} -algebra L is called *quasi-strictly right-sided (qsrs-algebra)* provided that $a \leq (\top_L \rightarrow_I a) \otimes \top_L$ for every $a \in L$.

Theorem 24

There exists a functor $L\text{-FC}_i^{\text{DMR}} \xrightarrow{H_{DC}^i} \mathbf{S}\text{-FC}^{\text{C}}$, which is given by $H_{DC}^i(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = (L^{G_1}, L^{M_1}, L, \hat{l}_1) \xrightarrow{(\alpha, \beta, 1_L)} (L^{G_2}, L^{M_2}, L, \hat{l}_2)$, where $\hat{l}_j(s, t) = \top_L$ if (s, t) is a formal concept of \mathcal{K}_j , and \perp_L otherwise. If L is a qsrs-algebra, then H_{DC}^i is a (non-full) embedding.

From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

Definition 23

- Given an \mathbf{L} -algebra L , $L\text{-FC}_i^{DMR}$ is a subcategory of $L\text{-FC}^{DMR}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have injective maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{op}} L^{M_1}$.
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From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

Definition 25

Given an \mathbf{L} -algebra L , $L\text{-FC}_{rfp}^{DMR}$ is a subcategory of $L\text{-FC}^{DMR}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{op}} L^{M_1}$ such that $K_2 \circ H_2 \circ \alpha(s) = \alpha(s)$ implies $K_1 \circ H_1(s) = s$, and $H_1 \circ K_1 \circ \beta^{op}(t) = \beta^{op}(t)$ implies $H_2 \circ K_2(t) = t$, for every $s \in L^{G_1}$, $t \in L^{M_2}$.

Theorem 26

There exists a functor $L\text{-FC}_{rfp}^{DMR} \xrightarrow{H_{DC}^{rfp}} \mathbf{S}\text{-FC}^C$, which is given by $H_{DC}^{rfp}(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = (L^{G_1}, L^{M_1}, L, \hat{h}_1) \xrightarrow{(\alpha, \beta, 1_L)} (L^{G_2}, L^{M_2}, L, \hat{h}_2)$, where $\hat{h}_j(s, t) = \top_L$ if (s, t) is a formal concept of \mathcal{K}_j , and \perp_L otherwise. If L is a qsr s-algebra, then the functor is a (non-full) embedding.

From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

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Given an \mathbf{L} -algebra L , $L\text{-FC}_{rfp}^{DMR}$ is a subcategory of $L\text{-FC}^{DMR}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{op}} L^{M_1}$ such that $K_2 \circ H_2 \circ \alpha(s) = \alpha(s)$ implies $K_1 \circ H_1(s) = s$, and $H_1 \circ K_1 \circ \beta^{op}(t) = \beta^{op}(t)$ implies $H_2 \circ K_2(t) = t$, for every $s \in L^{G_1}$, $t \in L^{M_2}$.

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From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

Definition 27

Given an \mathbf{L} -algebra L , $L\text{-FC}_{orp}^{DMR}$ is a subcategory of $L\text{-FC}^{DMR}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have order-preserving maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{op}} L^{M_1}$.

Theorem 28

There exists a functor $L\text{-FC}_{orp}^{DMR} \xrightarrow{H_{DC}^{orp}} \mathbf{S}\text{-FC}^C$, which is given by $H_{DC}^{orp}(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = (L^{G_1}, L^{M_1}, L, \hat{l}_1) \xrightarrow{(\alpha, \beta, 1_L)} (L^{G_2}, L^{M_2}, L, \hat{l}_2)$, where $\hat{l}_j(s, t) = \top_L$ if (s, t) is a formal preconcept of \mathcal{K}_j , and \perp_L otherwise. If L is a $qsrs$ -algebra, then the functor is a (non-full) embedding.

From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

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Given an \mathbf{L} -algebra L , $L\text{-FC}_{orp}^{DMR}$ is a subcategory of $L\text{-FC}^{DMR}$, with the same objects, and whose morphisms $\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2$ have order-preserving maps $L^{G_1} \xrightarrow{\alpha} L^{G_2}$, $L^{M_2} \xrightarrow{\beta^{op}} L^{M_1}$.

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From $\mathbf{S}\text{-FC}^{DMR}$ to $\mathbf{S}\text{-FC}^C$

Theorem 29

There exists a functor $L\text{-FC}^{DMR} \xrightarrow{H_{DC}} \mathbf{S}\text{-FC}^C$, which is given by $H_{DC}(\mathcal{K}_1 \xrightarrow{f} \mathcal{K}_2) = (L^{G_1}, L^{M_1}, L, \hat{l}_1) \xrightarrow{(\alpha, \beta, 1_L)} (L^{G_2}, L^{M_2}, L, \hat{l}_2)$, where $\hat{l}_j(s, t) = \top_L$ if (s, t) is a formal protoconcept of \mathcal{K}_j , and \perp_L otherwise. If L is a qsrs-algebra, then H_{DC} is a (non-full) embedding.

Final remarks

- This talk considered some approaches to morphisms of lattice-valued formal contexts of Formal Context Analysis (FCA).
- We constructed several categories, whose objects are lattice-valued analogues of formal contexts of FCA, and whose morphisms reflect the crisp setting of Chu spaces, the lattice-valued setting of J. T. Denniston, A. Melton, and S. E. Rodabaugh, as well as the many-valued setting of B. Ganter and R. Wille.
- We considered a number of functors between the categories of formal contexts, embedding each of them into its counterparts.

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FCA without relations

The difference between the settings of relations and Galois connections in the lattice-valued case, motivates the following problem.

Problem 30

Is it possible to build a lattice-valued approach to FCA, which is based in order-reversing Galois connections on lattice-valued powersets, which are not generated by lattice-valued relations on their respective sets of objects and their attributes?

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Thank you for your attention!